RESEARCH ARTICLE

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Optimum FIR Filtersfor Digital Pulse Compression of Biphase Barker Codes with Low Sidelobes

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Abstract

In Wireless signals and Radar signals where power, real estate, speed and low cost are tight constraints and Doppler tolerance is not a major concern biphase codes are popular and FIR filter is used for digital pulse compression (DPC) implementation to achieve required range resolution. Disadvantage of low peak to sidelobe ratio (PSR) of biphase codes can be overcome by linear programming for either single stage mismatched filter or two stage approach i.e. matched filter followed by sidelobe suppression filter (SSF) filter. Linear programming (LP) calls for longer filter lengths to obtain desirable PSR. Longer the filter length greater will be the number of multipliers, hence more will be the requirement of logic resources used in the FPGAs and many time becomes design challenge for system on chip (SoC) requirement.

This requirement of multipliers can be brought down by clustering the tap weights of the filter by kmeans clustering algorithm at the cost of few dB deterioration in PSR. The cluster centroid as tap weight reduces logic used in FPGA for FIR filters to a great extent by reducing number of weight multipliers. Since kmeans clustering is an iterative algorithm, centroid for weights cluster is different in different iterations and causes different clusters. This causes difference in clustering of weights and sometimes even it may happen that lesser number of multiplier and lesser length of filter provide better PSR.

In this paper few sample optimum biphase codes have been provided in tabular form with their optimum sidelobe supression filter (SSF) with optimum lengths and minimum multipliers to achieve low sidelobe level of -35 dBto -40 dB and have been compared against PSR achieved without clustering. Clustering is used in FIR filter for pulse compression and its effect is seen on peak to sidelobe ratio. This is being a genericmethod for FIR filters to reduce number of multipliers can be extended to other similar applications.

IndexTerms— Digital Pulse Compression (DPC), FIR filter, range sidelobes, sidelobe suppression filter(SSF), range sidelobe reduction, Kmeans algorithm, clustering, centroid, biphase codes, peak to sidelobe ratio (PSR).

I. I INTRODUCTION

Even though biphasecodes are used for digital pulsecompression in radars and wireless applications for quite some time and various methods have been suggested in literature to reduce their peak to sidelobe ratio, still designer struggles to achieve low peak to sidelobe ratio for biphase codes in the absence of readily available sidelobe suppression methodology filter design and its efficient implementation method. In applications where system is supposed to be built on a single chip in order to meet the specifications of real estate, power, cost and speed, each and every module of the system has to be efficient implementation wise. Efficient implementation of DPC has been suggested in [4] for 13 bit barker code but it is quite exhaustive to group the multipliers when SSF length increases and inefficient grouping of multiplier weights causes severe degradation in PSR, achieved with linear programming. So a method has been evolved to achieve optimal solution, taking care of all the constraints. Kmeans is an algorithm finds application

in image processing for clustering of data, has been used to provide solution to cluster the multipliers used as tap weights for efficient implementation and without much degradation in PSR.

Iterative Kmeans algorithm causes difference in clustering of weights in each iteration and sometimes even it may happen that lesser number of multiplier and lesser length offilter provide better PSR. Optimum length of SSF with minimum multipliers to achieve minimum PSR is desirable. However when all the multipliers will be implemented, there will be no deterioration in PSR achieved with linear programming provided overall design fits in the FPGA. This approach is useful in applications where System on Chip (SoC) via FPGA is supposed to be realized or in applications where design is being finalized for ASIC for bulk production, so logic required for each module of system should be minimum.

II. RANGE SIDELOBE SUPPRESSION FOR TWOSTAGE APPROACH

Range sidelobe suppression in DPC for biphase codes can be achieved by single stage mismatched filter or by two stage approach as shown in figure 1. Paper provides solution for two stage approach and optimal SSF filter design.

Multiplier weights for single stage filter can also be achieved in similar manner.

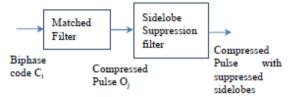


Figure (1) : Two stage approach

Sidelobe suppression filter is designed in two steps. In first step its multiplier weights are obtained by linear programming. In next step multiplier weights are grouped together by kmean clustering algorithm to bring down multipliers requirement. Multiplier weights obtained offline can be stored in memory to be used as shown in figure 2 for column 3 of table 2 with 6 weights.

III. DESIGN STEPS OF SSF

Linear programming

In step one linear programming technique of optimization is used to compute filter tap weights for binary phase coded waveforms.

Transmitted binary code,

 $Ci = \pm 1, l \leq i \leq N.$

Pulse compression filter output,

 $O_{j,l} \leq j \leq 2N-1$.

Sidelobe suppression filter tap weights,

 $S_{j,l} \leq j \leq M$, $M \geq 2N-1$.

Ci = 0 if i < 1 or i > N. Similarly, Oj = 0 if j < 1

orj > 2N - 1.

To obtain the tap weights for SSF constraints can be given as below.

Maximize
$$\sum_{J=1}^{M} S_{J} O_{j-(M-2N+1)/2}$$
 (1)

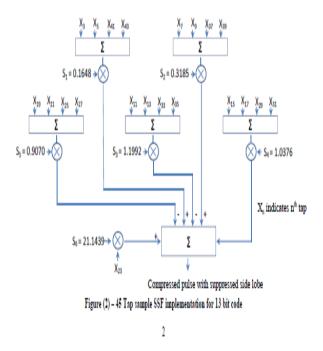
Subject to $\left|\sum_{i=1}^{M} S_{j} O_{j-k}\right| \leq p$

$$2(1 - N) \le k \le (M - 1), \ k \ne (M - 2N + 1)/2$$

Where p is the max sidelobe level for the code atcompression filter output. The objective function to bemaximized in equation (1) is the mainlobe response oftransmitted waveform centered in the filter weights.Equation 1 is having linear objective function with Mvariables and more than M linear inequality constraintsand can be solved by simplex algorithm of linearprogamming.

Kmeans algorithm

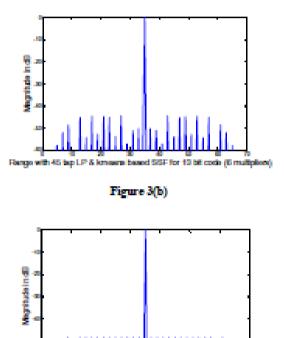
In step two k-means clustering [10] has been used to getcentroid to be used as tap weight. K-means is a two-phaseiterative algorithm to minimize the sum of point-to-centroid distances, summed over all K clusters. Thus, the first phaseis thought of as providing a fast but potentially onlyapproximate solution as a starting point for the secondphase. In second phase points are individually reassigned ifdoing so will reduce the sum of distances, and clustercentroids are recomputed after each reassignment.



To achieve approximately -40 dB PSR via optimum solution in terms of minimum SSF length and minimum multipliers exhaustive search is done for biphase codes mentioned in table 1 and few optimized results are tabulated in tables 2-4. Hardware implementation of sample SSF has been shown infigure 2.

Table 1 B	Biphase (Code used
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Code Length (Bit)	Biphase Code (hex)
13	1F35
25	1CE0549
51	71C077376ADB4
73[2]	1E78C0224975217C4E4
83[2]	4463BBF70F17E24FA4ADA
99[2]	63912B7FBDB06B6268FAA8D71



-												I		
	0		20	100							1			
90	45	tep	LP.		F	13	1	10	0 0	• (ore)	

Figure 3(a)

It can be noted from table 2, that for odd lengths of SSF, multiplier weights are symmetrical on both sides of middle tap weight while for even length of SSF, multiplier weights may or may not be symmetrical on both sides of middle tap weight. Hence, for 25 and 85 bit codes mentioned in table 1, SSF weights have been provide for odd SSF lengths up to middle tap weight in table 3.

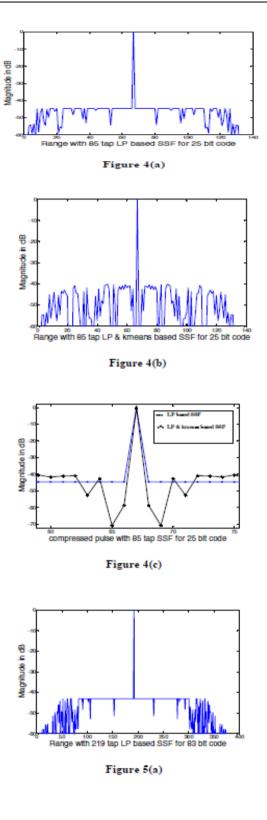
Index	LP b assed 45 tap SSF 12 Muhtipiliers PSR= 48.37 dB	LP & kmeans based 45 tap SSF 6 Multipli er PSR= - 44.52 dB	LP & kine and based 45 tap SSF 9 Multipli er PSR= - 45.19 dB	LP & kineans based 48 (ap SSF 8 Multipli er PSR= - 44.26 dB	LP & Kimeans based 44 TapSSF 8 mull fipl ier PSR=-43.40 dB	LP & kmeans based 42 tapSSF10 Multipli er PSR= - 46.83 dB	LP & Kineans based 42 TapSSF 7 multiplier PSRe-43.46 dB	LP based 37 Tap SSF 10 Multipli er PSR=-43.13 dB	LP & Kmeans based 37 Tap SSF 4 multiplier PSR=40.0 dB
1	0.0332	0.0000	0.0028	-0.0049	-0.0000	0.0885	0.1291	0.2222	0.2811
23	-0.0000 0.1235	0.0000 0.1648	-0.0028	0.0049 0.0049	0.1371 -0.0000	0.0000	0.0000 0.1291	-0.0000	0
	0.1235	0.1648	0.1235	0.0049	-0.0000	0.1696	0.1291	0.2724	0.2811
4	-0.0000 0.2061	0.0000 0.1648	-0.0028 0.2430	0.1694 0.0049	0.1371	-0.0000 0.2690	-0.0000 0.3709	0.0000 0.3164	0
6	-0.0000	0.0000	-0.0028	0.1694	0.2509	0.0000	0.0000	0.0000	0.2811
17	0.2798	0.3185	0.2430	-0.0049	-0.0000	0.4729	0.3709	-0.6498	-1.1609
8	-0.0000	0.0000	-0.0028	0.3298	0.4821	-0.0000	-0.0000	0.0000	0
9	0.3573	0.3185	0.3573	0.0049	-0.0000	-1.0470	-1.0089	-0.6057	-1.1609
10	-0.0000	0.0000	-0.0028	0.3298	-1.0345	-0.0000	-0.0000	-0.0000	0
11 12	-1.2380	-1.1992	-1.1992	0.0049	-0.0000	-0.9708 -0.0000	-1.0089 -0.0000	-0.5556	-1.1609
13	0.0000	0.0000	0.0028	0.0049	-0.9283	-0.0000	-1.0089	-0.0000 -0.6667	0 -0.9351
14	-0.0000	0.0000	-0.0028	-1.1674	-1.0345	-0.0000	-0.0000	0.0000	0.9551
15	-1.0841	-1.0376	-1.0841	-0.0049	-0.0000	-0.9708	-1.0089	-0.6019	-0.9351
16	-0.0000	0.0000	-0.0028	-1.0524	-1.0345	0.0000	0.0000	0.0000	0
17	-0.9912	-1.0376	-0.9607	-0.0049	0.0000	-0.8709	-0.8709	-0.5287	-0.9351
18	-0.0000	0.0000	-0.0028 -0.8837	-1.0524 -0.0049	-0.9283	0.0000	0.0000	0.0000 11.5944	0
19	-0.8837 0.0000	-0.9070 0.0000	-0.8837	-0.0049	0.0000	-0.7679 0.0000	0.0000	-0.0000	21.1439 0
20 21	-0.9303	-0.9070	-0.9607	-0.0049	0.0000	18.0584	18.0584	-0.5287	-0.9351
22	0.0000	0.0000	0.0028	-0.8521	18.1773	0.0000	0.0000	-0.0000	0
23	21.1439	21.1439	21.1439	-0.0049	0.0000	-0.7679	-0.7679	-0.6019	-0.9351
24	0.0000	0.0000	0.0028	21.2314	-0.7807	0.0000	0.0000	-0.0000	0
25 26	-0.9303 0.0000	-0.9070	-0.9607 0.0028	0.0049	-0.0000	-0.8709 -0.0000	-0.8709 -0.0000	-0.6667	-0.9351
20	-0.8837	0.0000	-0.8837	0.0049	-0.7807 -0.0000	-0.0000	-0.0000	0.0000	-1.1609
28	-0.0000	0.0000	-0.0028	-1.0524	-0.7807	0.0000	0.0000	0.0000	0
29	-0.9912	-1.0376	-0.9607	0.0049	0.0000	-1.0470	-1.0089	-0.6057	-1.1609
30	0.0000	0.0000	0.0028	-1.1674	-0.9283	-0.0000	-0.0000	-0.0000	0
31	-1.0841	-1.0376	-1.0841	0.0049	-0.0000	-0.9708	-1.0089	-0.6498	-1.1609
32 33	-0.0000 -1.1605	0.0000	-0.0028 -1.1992	-1.0324 0.0049	-1.0345	-0.0000 -1.0470	-0.0000 -1.0089	-0.0000 0.3164	0
34	-0.0000	0.0000	-0.0028	-1.1674	-1.0345	-0.0000	-0.0000	0.0000	0.2811
35	-1.2380	-1.1992	-1.1992	0.0049	-0.0000	0.4729	0.3709	0.2724	0.2811
36	-0.0000	0.0000	-0.0028	-1.2383	0.2509	-0.0000	-0.0000	0.0000	0
37	0.3573	0.3185	0.3573	0.0049	0.0000	0.2690	0.3709	0.2222	0.2811
38	-0.0000	0.0000	-0.0028	0.3298	0.2509	0.0000	0.0000		
39 40	0.2798	0.3185	0.2430	0.0049 0.3298	0.0000	0.1696	0.1291 0.0000		
40	-0.0000 0.2061	0.0000 0.1648	0.2430	0.3298	0.1371	0.0000	0.0000		
42	-0.0000	0.0000	-0.0028	0.1694	0.1371	0.0000	0.0000		
43	0.1235	0.1648	0.1235	-0.0049	-0.0000				
44	-0.0000	0.0000	-0.0028	0.1694	0.1371				
45	0.0332	0.0000	0.0028	-0.0049					
46 47				0.0049					
48				-0.0049					
10				0.0012					

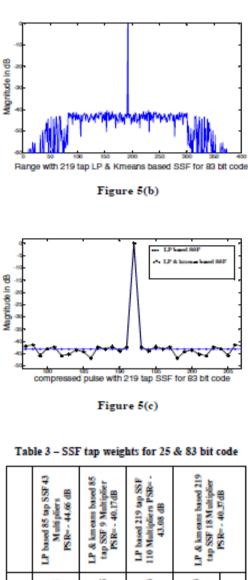
Table 2 - Multiplier weights for 13 bit binhase code

IV. SIMULATION RESULTS

2

Figure 4(a) and 4(b) shows range vs magnitude plot for 25 bit biphase code for 85 tap SSF for LP and LP &kmeans based SSF respectively, figure 5(a) and 5(b) shows range vs magnitude plot for 83 bit biphase code for 219 tap SSF for LP and LP &kmeans based SSF respectively. Figure 4(c) shows compressed pulse widths at the output of LP vs LP &kmeans based 85 tap SSF for 25 bit code, it can be seen thatcompressed pulse width doesn't change with LP andkmeans based SSF. Figure 5(c) shows compressed pulse at the output of LP vs LP and kmeans based 219 tap SSF for 83 bit code, same can be noted from figure 5(c) too that compressed pulse width remains unchanged with LP &kmeans based SSF.





	LP based 85 tap SSF 43 Multipliers PSR=- 44.66 dB	LP & km ears based 85 tap SSF 9 Muhiplier PSR=-40.17dB	LP based 219 tap SSF 110 Muhtipliers PSR= - 43.08 dB	LP & km eans based 219 (ap SSF 18 Multiplier	15.K= - 40.37dB
Index	Muhi plier Weight for 25 bit code	Multiplier Weight for 25 bit code	Multiplier Weight for 83 bit code	Multiplier Weight for 83 bit code	Multiplier Weight Index
1	0.2960	0.3213	0.0645	0.0520	11
12345678	0.1255	0.1169	-0.0307	-0.0134	3
3	-0.3293	-0.3213	-0.0865	-0.0913	4
4	0.0507	0.1169	-0.0004	-0.0134	3
5	0.5816	0.5218	-0.0464	-0.0520	11
6	-0.3629	-0.3213	0.0533	0.0520	11
7	0.0950	0.1169	-0.0972	-0.0913	4
8	0.2584	0.3213	0.1798	0.1818	5
9	-0.1135	-0.1169	-0.0207	-0.0134	3 14 11
10 11	0.2623	0.3213	0.1199	0.1323	14
11	0.1917	0.1169	-0.0572 0.0686	-0.0520	4
12 13	-0.2020 0.4804	0.5218		0.0913	1.7
14	0.1455	0.1169	-0.0153 0.0235	0.0134	3
15	-0.1053	-0.1169	-0.0738	-0.0913	4
16	0.1901	0.1169	0.0950	0.0913	4
17	0.1444	0.1169	-0.0403	-0.0520	n i
17 18	0.0093	0.1169	0.0672	0.0520	11

		, .	,	(//
19	-0.0890	-0.1169	-0.0494	0.0520	11
20	0.0395	0.1169	0.0966	0.0913	4
21	-0.1115	-0.1169	-0.0586	-0.0520	11
22	1.9624	1.9267	-0.0034	-0.0134	3
23	0.3754	0.3213	0.1173	0.1323	14
24	-1.4612	-1.4163	-0.0109	-0.0134	3
25	1.1472	1.2054	0.0470	0.0520	11
26	1.3275	1.4163	0.0573	0.0520	11
27	0.8788	-0.8134	0.1506	0.1323	14
28	1.4602	1.4163	0.0487	0.0520	11
29	-0.3218	-0.3213	-0.2787	0.2705	15
30	0.7988	0.8134	0.3156	0.3072	7
31	1.2636	1.2054	-0.2352	-0.2156	16
32	-0.1173	-0.1169	0.1382	0.1323	14
33	0.1395	0.1169	0.0397	0.0520	11
34	0.8471	0.8134	0.3109	0.3072	7
35	0.3644	0 3213	-0.4128	-0 1988	18
36	-1.0473	-1.0005	0.3945	0.3988	18
37	0.9865	1.0005	0.1236	0.1323	14
38	-0.9678	-1.0005	0.0786	0.0913	4
39	0.4433	0.5218	-0.4281	-0.4523	2
40	1.8910	1.9267	0 3619	0.3557	13
41	0.7289	-0.8134	0.2020	0.2156	16
42	-0.5819	-0.5218	0.2996	0.3072	7
43	15 9606	15 9606	-0 3696	-0 3557	13
44	A.C. 2000	10.0000	-0.0583	-0.0520	ñ
45			0.0894	0.0913	4
46			-0.2029	-0.2156	16
47			0.0364	0.0520	ĩĩ
48			-0 3533	-0 3557	13
49			0.2219	0.2156	16
50			-0.3386	-0.3557	13
- <u>5</u>			-0.2074	-0.2156	16
52			0 4048	0 3988	18
53			0.0436	0.0520	ii –
- 54			0.1277	-0.1323	14
55			0.1363	0.1323	14
56			0.0494	0.0520	11
57			-0 3893	-0 3988	18
58			-0.0694	-0.0520	ñ
59			0.1921	0.1818	5
60			-0 3966	-0.3988	18
61			-0.2695	-0.2705	15
62			0.2352	0.2156	16
63			-0.4760	-0.4523	2
64			0.0796	0.0913	4
65			-0.1921	-0.1818	5
66			0 3026	0.3072	ź
67			-0.1064	-0.0913	4
68			0.3889	0.0915	18
69			0.3412	0.3557	13
70			-0.2097	-0.2156	16
71			-0.1342	-0.1323	14
72			-1.0385	-1.0135	9
73			-0.0042	-0.0135	3
74			0.3635	0.3557	13
75			0.3655	0.2705	15
76			-0.1528	-0.1323	14
77			0.0827	0.0913	4
2 A			ALCONTRACTOR	WW0113	

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78		0.3852	0.3988	18
79		-1.0214	-1.0135	9
80		-0.5985	-0.6077	8
81		0.4217	0.3968	18
82		0.5931	0.6077	8
83		-0.5299	-0.5392	1
84		0.2627	0.2705	15
85		-0.3950	-0.3988	18
86		-0.0554	-0.0520	11
87		-0.8298	-0.8298	10
88		0.9804	1.0135	9 5
89		0.1869	0.1818	5
90		0.1016	0.0913	4
91		-0.1407	-0.1323	14
92		-0.0609	-0.0520	11
93		-0.1143	-0.1323	14
94		0.0112	0.0134	3
95		-0.1636	-0.1818	5
96		-0.0377	-0.0520	11
97		0.1092	0.0913	4
98		0.4520	0.4523	25
99		-0.1761	-0.1818	5
100		-0.0793	-0.0913	4
101		0.1102	0.0913	4
102		-0.0867	-0.0913	4
103		0.7487	0.7487	17
104		0.6314	0.6077	8
105		-0.3616	-0.3557	13
106		0.4532	0.4523	2
107		-0.2574	-0.2705	15
108		-0.2105	-0.2156	16
109		-0.5485	-0.5392	1
110		11.8514	11.8514	6

Table 4 - 13, 25, 51, 73, 83, 99 bit code results

Code length (bit)	tap SSF length	With LP PSR (dB)	Multi- pliers required	With LP & Kine and PSR (dB)	Mal6- pliers required
13	37	43.13	10	-43.06	10
	38	43.13	10	-42.93	9
	41	46.98	11	-41.87	9
	41	46.98	11	-42.15	6
	41	-46.98	11	-40.91	5
	45	48.37	12	-44.52	7
	45	-48.37	12	-42.82	6
	49	48.63	13	-45.16	8
25	85	-44.66	43	-40.17	9
	87	-46.38	44	-42.26	12
	89	4651	45	-44.24	23
	90	4651	90	-43.17	19
	91	4651	46	-41.66	16
	93	-46.88	47	-45.00	22
	97	-47.03	49	-40.32	11
	97	47.03	49	-45.19	25

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Jineshwar Nyamagoud Int. Journal of Engineering Research and Applications
ISSN : 2248-9622, Vol. 4, Issue 11(Version 1), November 2014, pp.102-108

51	175	-41.95	88	-41.01	39
	177	-42.68	89	-40.00	28
	179	-42.76	90	-41.27	38
	183	-42.93	92	-40.58	26
	183	-42.93	92	-41.38	36
	199	-43.44	100	-40.13	23
	199	-43,44	100	-41.88	39
	199	-43,44	100	-42.49	46
	201	-43.48	101	-41.15	29
73	223	-41.33	112	-40.33	52
	237	-42.21	119	-40.10	-30
	264	-43.32	264	-40.63	41
	269	-43.46	135	-40.02	-59
	272	-43.71	272	-40.21	36
	273	-43.91	137	-40.10	23
	273	-43.91	137	-40.85	31
	273	-43.91	137	-41.36	45
	273	-43.91	137	-42.43	-57
	274	-43.91	274	-40.27	25
	274	-43.91	274	-41.64	-54
	275	-43.92	138	-40.00	22
	275	-43.92	138	-41.14	37
	275	-43.92	138	-42.04	55
	281	-44.31	141	-40.25	37
	281	-44.31	141	-41.43	52
	281	-44.31	141	-42.00	60
	281	-44.31	141	-43.03	64
	289	-44.57	145	-40.27	40
83	177	-40.01	89	-39.50	38
	205	-41.36	103	-40.02	31
	207	-42.26	104	-40.30	25
	219	-43.08	110	-40.37	18
	255	-44.43	128	-40.40	15
	287	-45.82	144	-40.76	13
	211 216	42.39	106	-41.03	33
	210	-42.77 -43.16	216 112	-41.48 -41.18	34 25
	225	-43.10		-41.18	
			254		19
	288 228	-45.82 -43.19	288 228	-41.12 -42.17	16 37
	228	-43.19	126	42.48	27
	295	-46.39	148	-42.93	23
	325	-47.18	163	-42.90	20
<u> </u>	325	-47.19	326	-42.90	15
	251	-44.34	126	-43.08	38
	279	-45.10	140	-43.00	33
	295	-46.39	148	-43.04	32
	330	47.23	330	-43.71	24
	287	-45.82	144	-44.59	46
	306	-46.62	306	-44.00	38
	330	-47.23	330	-44.39	30
	295	-46.39	148	-45.30	63
	311	-46.65	156	-45.00	57
	321	-47.09	161	-45.01	45
99	366	-40.11	366	-36.41	79
	367	-40.38	184	-38.31	76
		19.39			- * *

381	-41.18	191	-39.07	70
381	-41.18	191	-39.64	84
381	41.18	191	-39.14	65
419	41.96	210	-40.02	104
423	-42.00	212	-39.11	79
429	-42.07	215	-40.22	105
437	-42.16	219	-40.04	88
437	-42.16	219	-40.69	108
439	42.24	220	-40.51	95
475	-42.87	238	-40.06	95
475	42.87	238	-40.47	106
491	43.18	246	-40.68	118

V. CONCLUSION

Biphase and polyphase codes will never be out of fashion in DPC, irrespective of radar history as they are foundation stone of DPC. In applications where Doppler tolerance is not a major concern biphase codes are preferred over polyphase codes for their simplicity. Fast Convolution is preferred over FIR filters for digital pulse compression in radars for computational advantage for longer codes and range gates in digital signal processing. However for FPGA based implementation of matched filter along with SSF for longer biphase codes, poses constraint over logic resources used. The proposed optimization solution paves a way to reduce hardware in the generation of radar where multiple current beams simultaneous are formed. Multiple simultaneous beams in radars play important role, especiallyin difficult to detect and multi target environment. This method of reducing number of multiplier is applicable to polyphase codes too. But polyphase codes do not have choice to improve PSR using longer length SSF; code length and incorporation of window at receiver determines PSR. Achieved PSR deteriorates a bit because of clustering.

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